

TYPES 2022

Partial Dijkstra monads for all



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How to define effects in F^* ?



How to use effects in F*?

```
using state = nat
```



```
let decr () : unit  
= let x = get () in  
  put (x - 1)
```

How to use effects in F*?

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using state = nat
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```
let decr () : St unit  
= let x = get () in  
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How to use effects in F*?

```
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```
let decr () :  
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  (requires  $\lambda s_0. 1 \leq s_0$ )  
  (ensures  $\lambda s_0 s_1. s_1 = s_0 - 1$ )  
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```

```
get () : St nat (requires  $\lambda s_0. \top$ )  
        (ensures  $\lambda s_0 s_1 r. s_1 = s_0 = r$ )
```

```
put (x : nat) : St unit (requires  $\lambda s_0. \top$ )  
                (ensures  $\lambda s_0 s_1 r. s_1 = x$ )
```



```
bind (get ()) (λ x,  
  put (x - 1)  
)
```



```
let decr () :  
  St unit  
  (requires λ s0. 1 ≤ s0)  
  (ensures λ s0 s1 r. s1  
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  put (x - 1))
```

```
get () : St nat (requires λ s0. T)  
  (ensures λ s0 s1 r. s1 = s0 = r)
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put (x : nat) : St unit (requires λ s0. T)  
  (ensures λ s0 s1 r. s1 = x)
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bind (get ()) (λ x,  
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put (x : nat) : St unit (requires λ s0. T)  
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```

Computation level

```
bind (get ()) (λ x,  
  put (x - 1)  
)
```

Specification level

bind

(requires λ s₀. T)

(ensures λ s₀ s₁ r. s₁ = s₀ = r)

(requires λ s₀. T)

(ensures λ s₀ s₁ r. s₁ = x)

(requires λ s₀. 1 ≤ s₀)

(ensures λ s₀ s₁ r. s₁ = s₀ - 1)

Computation level

```
bind (get ()) (λ x,  
  put (x - 1)  
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Specification level

bind

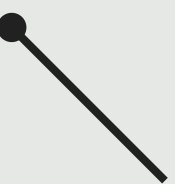
..... \equiv $(\text{requires } \lambda s_0. 1 \leq s_0)$
 $(\text{ensures } \lambda s_0 s_1 r. s_1 = s_0 - 1)$

$(\text{requires } \lambda s_0. \top)$
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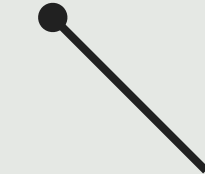
Dijkstra monads

$D (A : \text{Type}) (w : W A) : \text{Type}$


ordered specification *monad*

Dijkstra monads

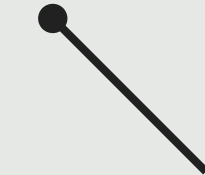
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 ordered specification *monad*

$\text{return}^D (x : A) : D A (\text{return}^W x)$

Dijkstra monads

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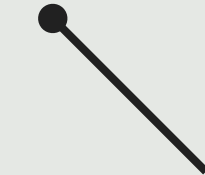
 ordered specification *monad*

$\text{return}^D (x : A) : D A (\text{return}^W x)$

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Dijkstra monads

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 ordered specification *monad*

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$\text{subcomp}^D (c : D A w) : \text{Pure } (D A w')$
 $(\text{requires } w \leq w')$
 $(\text{ensures } \lambda r, r = c)$

Dijkstra monads *for all*

(Maillard et al. ICFP 2019)

$$M \xrightarrow{\theta} W$$

$$D (A : \text{Type}) (w : W A) = \{ c : M A \mid \theta c \leq w \}$$

e.g.

computation monad

St

specification monad

$$M A = S \rightarrow A \times S$$

$$\text{return}^M x = \lambda s_0. (x, s_0)$$

$$\text{bind}^M c f = \lambda s_0.$$

$$\text{let } (x, s_1) = c s_0 \text{ in } f x s_1$$

$$W A = \text{post } A \rightarrow \text{pre}$$

$$\text{return}^W x = \lambda P s_0. P x s_0$$

$$\text{bind}^W w wf =$$

$$\lambda P. w (\lambda s_1 x. wf x P s_1)$$

Dijkstra monads *for all*

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$$\theta c = \lambda P s_0. \text{let } (x, s_1) = c s_0 \text{ in } P x s_1$$

```
using state = nat
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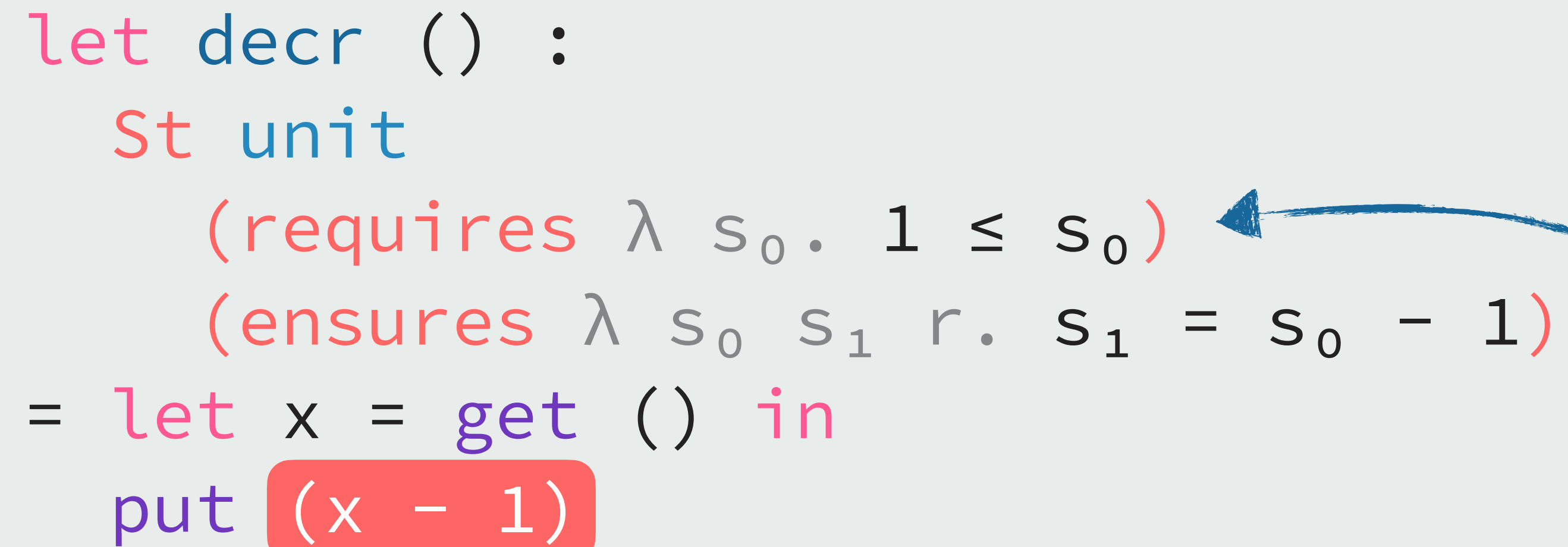
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  St unit  
  (requires  $\lambda s_0. 1 \leq s_0$ )  
  (ensures  $\lambda s_0 s_1 r. s_1 = s_0 - 1$ )  
= let x = get () in  
  put (x - 1)
```

```
using state = nat = { x : int | 0 ≤ x }
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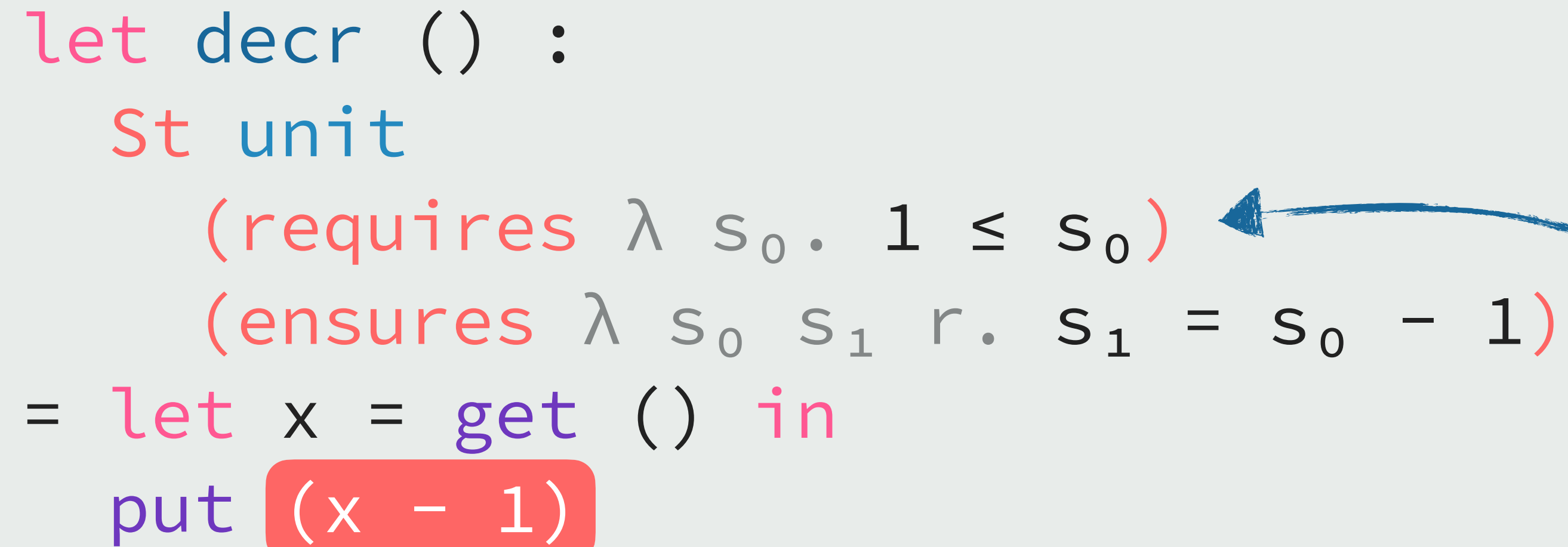
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ill-defined without pre-condition

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ill-defined without pre-condition

```
let head (l : list A) :  
  Pure A (requires l ≠ []) (ensures λ r. T)  
= match l with  
| x :: l' → x
```

```
let decr () :  
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let head (l : list A) :  
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```
let x = get () in  
let y = get () in  
assert (x = y)
```

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let decr () :  
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```
let x = get () in  
let y = get () in  
assert (x = y) ← need the pre-condition deeply
```

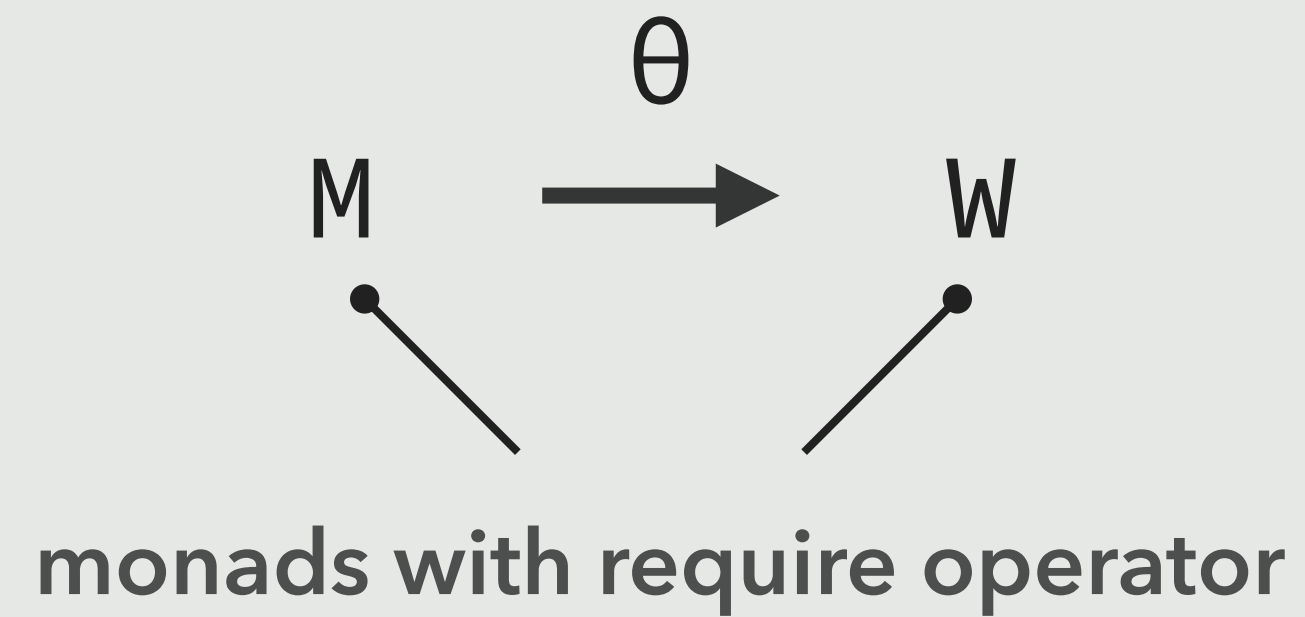


Dijkstra monads *for all*

$$M \xrightarrow{\theta} W$$

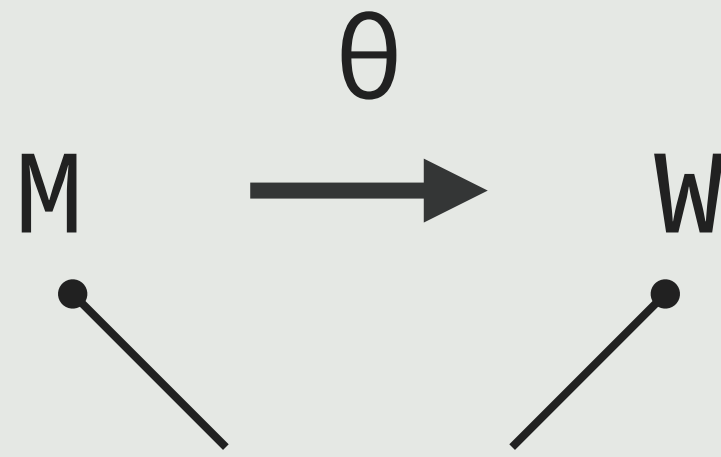
$$D (A : \text{Type}) (w : W A) = \{ c : M A \mid \theta c \leq w \}$$

Partial Dijkstra monads *for all*



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Partial Dijkstra monads *for all*

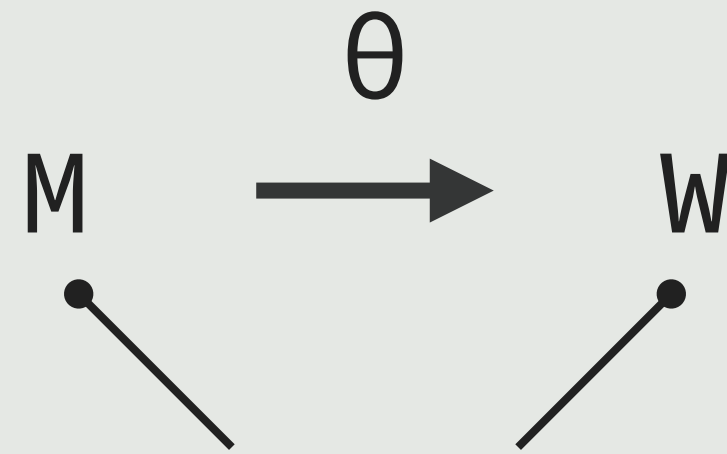


monads with require operator

`requireM (P : \mathbb{P}) : M P`

`D (A : Type) (w : W A) = { c : M A | $\theta c \leq w$ }`

Partial Dijkstra monads *for all*



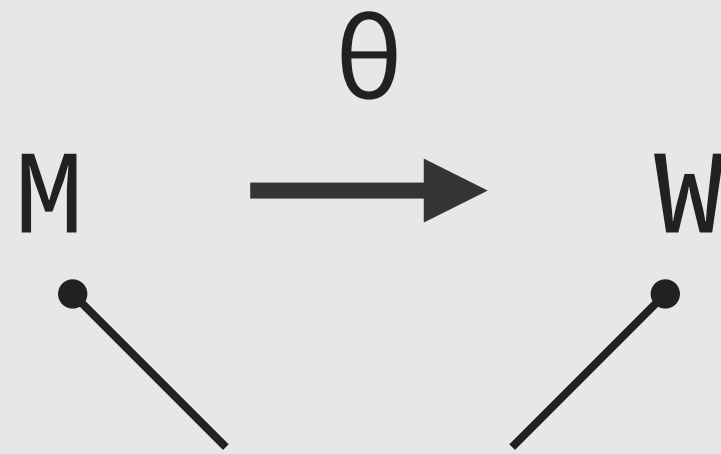
monads with require operator

$\text{require}^M (P : \mathbb{P}) : M P$

$\text{require}^D (P : \mathbb{P}) : D P (\text{require}^W P)$

$D (A : \text{Type}) (w : W A) = \{ c : M A \mid \theta c \leq w \}$

Partial Dijkstra monads *for all*



monads with require operator

$\text{require}^M (P : \mathbb{P}) : M P$

$\text{require}^D (P : \mathbb{P}) : D P (\text{require}^W P)$

$\theta (\text{require}^M P) = \text{require}^W P$

$D (A : \text{Type}) (w : W A) = \{ c : M A \mid \theta c \leq w \}$

State partial Dijkstra monad

$$M \ A = S \rightarrow G \ (A \times S)$$

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$$M A = S \rightarrow G (A \times S)$$

where $G B = \Sigma (\text{pre} : \mathbb{P}). \text{pre} \rightarrow B$

State partial Dijkstra monad

$$M A = S \rightarrow G (A \times S)$$

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given by (pre, f) with $f : \text{pre} \rightarrow B$

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$$\theta c = \lambda P s_0,$$

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$\text{let } (\text{pre}, f) := c s_0 \text{ in}$

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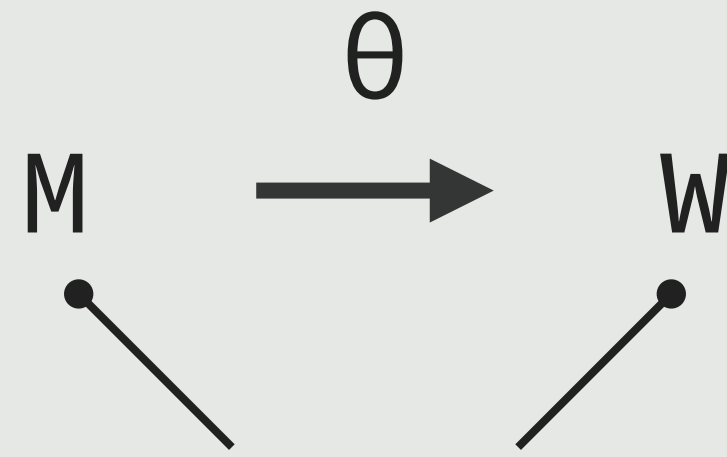
$\theta c =$

$\lambda P s_0,$

$\text{let } (\text{pre}, f) := c s_0 \text{ in}$

$\exists (h : \text{pre}). \text{let } (x, s_1) := f h \text{ in } P x s_1$

Partial Dijkstra monads *for all*

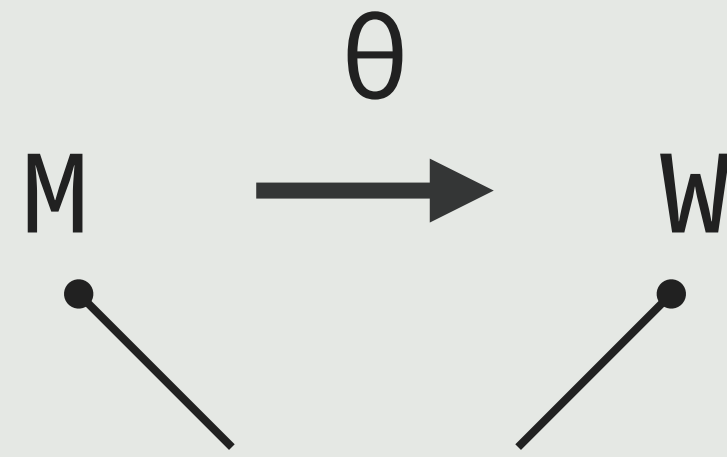


monads with require operator



partial effect in F^*

Partial Dijkstra monads *for all*

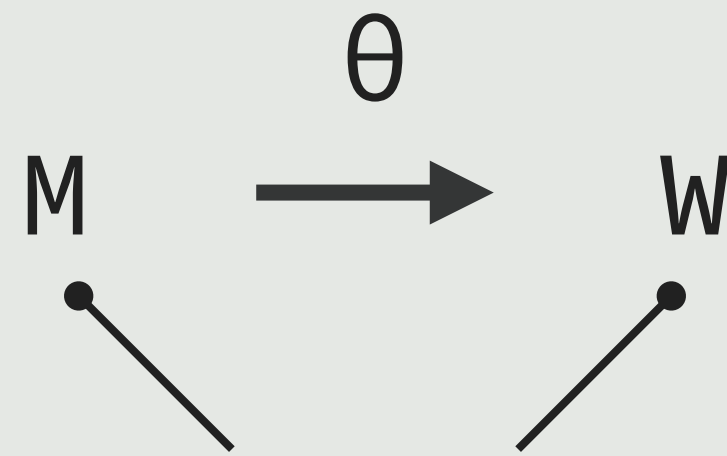


monads with require operator



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Partial Dijkstra monads *for all*



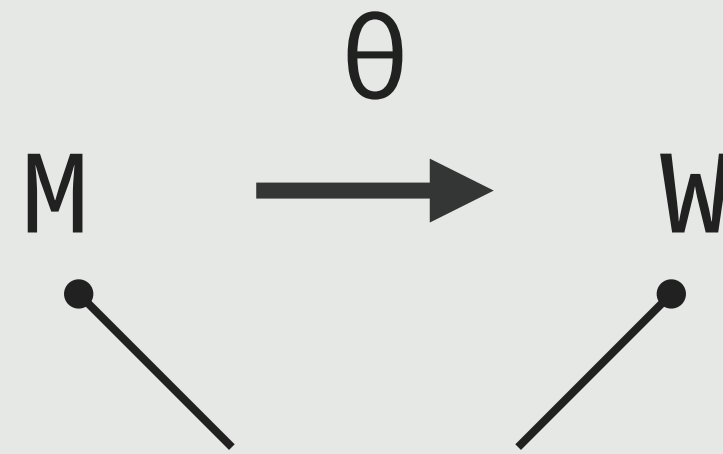
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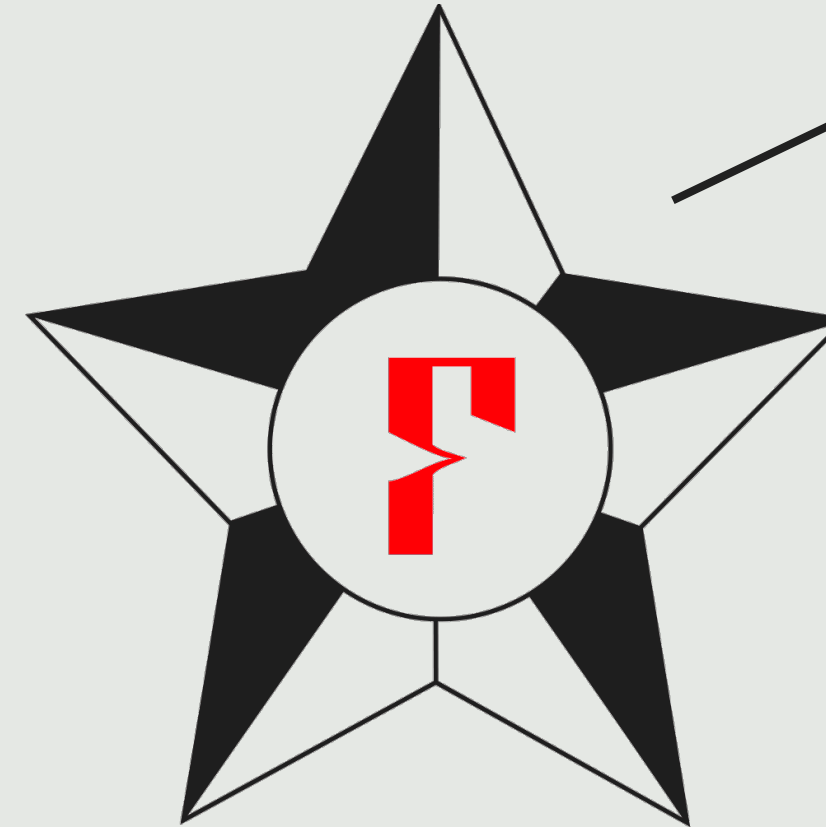
partial effect in F^*

(Ahman *et al.* POPL 2017)

Partial Dijkstra monads *for all*



monads with require operator



partial effect in F^*

DM4Free

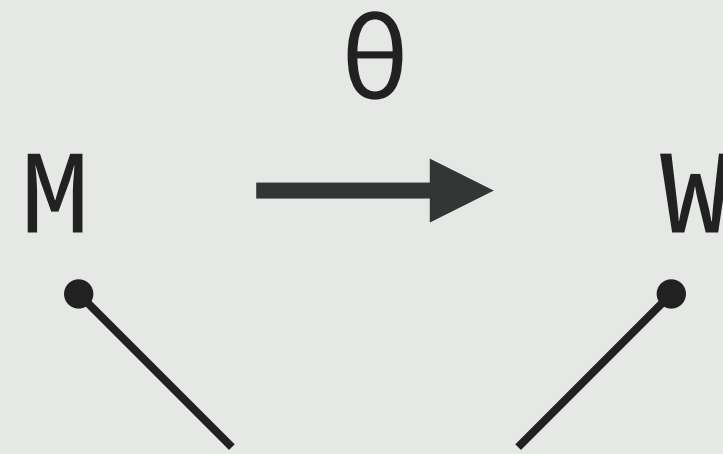
State, exceptions...

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Free monads

IO, non-determinism...

Partial Dijkstra monads *for all*

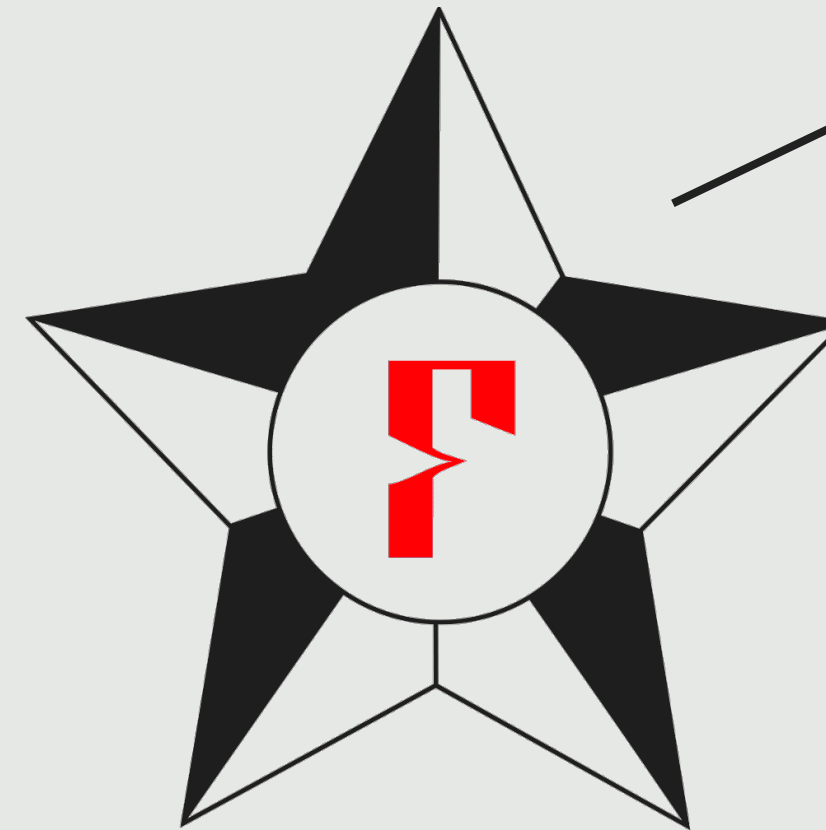


monads with require operator



DM4Free
State, exceptions...

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partial effect in F^*

Free monads
IO, non-determinism...

Non-terminating IO